

Institut Mines-Téléco



Pallas: HPC Trace analysis at scale

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Context: performance analysis at exascale Investigating performance problems

- Many sources of performance problems at scale
 - Collective communication, load imbalance, network contention, NUMA effects, ...
- Finding a bottleneck requires investigation
 - Run the application once, analyze its execution trace
- Applications mix programming models (MPI+OpenMP, MPI+CUDA), and libraries (StarPU, netcdf, Pytorch, ...)
- Need for generic, modulable tracing tools

Proposal: Pallas trace format

Detecting sequences of events for

Example of a Example of a sequential trace Pallas trace



MPI_Recv()

Analyzing huge performance data Finding the needle in the haystack

- Traces are records of execution made for post-mortem
 analysis
 - Store detailed information.
 - Can take huge amounts of storing space
 - Require a lot of processing before being useful
- Gets worse with execution time and number of threads used.
- Traces of HPC applications have recurring patterns, which means they have a lot of redundant data.

Analysis-optimized trace storage

Separating data from

1	Ч	nb_threads, compression, MPI_rank,	Archive head
		Tokens 🗾 🔲 ···	





better data analysis

- Intercepted functions are grouped by call signature as Tokens.
- Their **durations** are stored for analysis.
- Repetition of similar functions are detected using memcmp and stored as loops, allowing for better view of the structure of the trace.

I_Send()	200 * Sequence 1	
I_Recv()	MPI_Send()	
ompute()	MPI_Recv()	
I_Send()	<pre>compute()</pre>	
I_Recv()	1 * Sequence 2	
ompute()	MPI_Barrier()	
× 200	MPI_Recv()	
_Barrier()		

metadata

- One folder per MPI rank
- Header file with general information (lightweight)
- Data file with durations:
 - Efficient compression
 - Easily removable
 - On-demand loading



Evaluations

Comparison with **Pilgrim** (pattern recognition and compression) and **OTF2** (sequential format with timestamp encoding).

- Negligible overhead for most applications (1-10%) except for irregular ones (57% for Quicksilver)
- Trace size lies between OTF2 and Pilgrim
- Pilgrim intercepts less function and compresses all timestamps at once \rightarrow Smaller traces, better compression
- Compression can reduce Pallas trace size by a factor of 10

Trace analysis

We have implemented two analysis tools for OTF2, Pallas and Pilgrim: one to **plot a communication matrix**, the other to **detect contention**.

- Both Pilgrim and Pallas run the tools nearinstantaneously.
- Pallas uses a lot less memory than Pilgrim.













Conclusion

- Low overhead (<5%) for regular applications
- Reasonable trace size compared to alternatives
- Quick and memory-light analyses

Ongoing work

- Provide trace visualization tools
- Pattern detection for irregular applications
- Develop complex trace analysis methods
- Tracing non-MPI kernels (CUDA, StarPU etc)
- Performance analysis at scale (TBs of data)
- Provide better compression techniques

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